## Assignment 1 Grading

Grading Scheme 9 Qs, equally weighted, 2 pts per Q, total (18).

LaTeX comments Equations should end with full stops, e.g. let X be a scheme such that

$$H^{i}(X, \mathcal{N}_{X}) = 0$$
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L2, Ex 4 What I meanthere is not <u>necessarily</u> surjective, and I'm glad many noticed this.

L2, Ex5 What is up with all the people uniting X — Y? This is a strange (and in my opinion, bad) naming style, as most people will read xy as x o y. If you insist, while e.g. Sxy. But everybody just was Greek or Roman letters and so should you.

Yoneda uniqueness A few people got this, which was good. Let  $F: \mathcal{C} \xrightarrow{p} Set$ be a functor,  $C \in ob(\mathcal{C})$  and

## $\underline{\Phi}_{C,F}: \operatorname{Nat}(hc, F) \longrightarrow F(C), \quad \underline{\Phi}_{C,F}(\alpha) = \alpha_{C}(\operatorname{id} c)$

the usual Yoneda isomorphism. A natural (in C and F) family of bijections  $\mathcal{K}_{c,F} \subseteq \operatorname{Nat}(\operatorname{hc}, F)$  will give, upon precomposition, a second natural family of bijections  $\mathfrak{F} \circ \mathfrak{N}$ . So we want to know

 $\bigcirc$ 

It's havd to work on F because it is generic. But suppose we had a natural iso  $\mathcal{I}^{C}: h_{c} \longrightarrow h_{c}$  for each  $C \in ob(\mathcal{B})$  such that for every  $f: C \longrightarrow C'$  the diagram  $h_c \xrightarrow{\gamma} h_c$ hf | hf (2.1) hc' ----- hc' commutes. Then the family of bijections  $\mathcal{T}_{c,F}: \operatorname{Nat}(h_{c},F) \longrightarrow \operatorname{Nat}(h_{c},F), \quad \mathcal{T}_{c,F}(\alpha) = \alpha \cdot \mathcal{Z}^{C}$ is certainly natural in C, F and  $\mathcal{T}_{c,F}(\alpha)_c(id_c) = (\alpha \circ \gamma)_c(id_c) = \alpha(\gamma_c^c(id_c))$ so as long as  $\mathcal{N}_{\epsilon}^{\epsilon}(idc) \neq idc$  we have a good shot at (2) as well. So let wundertand what the data of { 2 c Nat (hc, hc) ) c c b (o) plus naturality (2.1) means. We know Nat(hc, hc)  $\cong$  hc(c)  $\cong$  Home(C,C) sending ? to  $Y_c = 2c(idc)$ , and  $2^c$  a natural iso  $\iff Y_c$  is an iso. Combining this with (2.1), wehave

(2)

Lemma There is a bijection between families of natural isos

 $\{\mathcal{Y}^{c}: h_{c} \rightarrow h_{c}\}_{c \in \delta}(\delta), \text{ natural in } C$ 



(The monoid Nat(ide, ide) is called the center of C. It is a commutative monoid ZC and the units we denote  $Z(C)^{\times}$ ).

Def Given a family of natural isos  $\{Y_c: c \rightarrow c\}_{c \in blo}$  as above, define

 $\underline{\oplus}_{c,F}^{\mathcal{Y}}: \operatorname{Nat}(h_{c},F) \longrightarrow F(C), \quad \underline{\oplus}_{c,F}^{\mathcal{Y}}(\alpha) := \alpha_{c}(\mathcal{Y}_{c}).$ 

This is a natural bijection, by construction, and

$$\begin{array}{rcl} \forall \mathsf{C}\forall\mathsf{F}\ \bar{\Phi}_{\mathsf{C},\mathsf{F}}^{\mathsf{Y}} = \bar{\Phi}_{\mathsf{C},\mathsf{F}} & \Longrightarrow & \forall \mathsf{C}\ \bar{\Phi}_{\mathsf{C},\mathsf{hc}}^{\mathsf{Y}} = \bar{\Phi}_{\mathsf{C},\mathsf{hc}} \\ & \Longrightarrow & \forall \mathsf{C}\ \bar{\Phi}_{\mathsf{C},\mathsf{hc}}^{\mathsf{Y}}(\mathsf{id}_{\mathsf{hc}}) = \bar{\Phi}_{\mathsf{C},\mathsf{hc}}(\mathsf{id}_{\mathsf{hc}}) \\ & \Longrightarrow & \forall \mathsf{C}\ \bar{\Psi}_{\mathsf{C}}^{\mathsf{Z}} = \mathsf{id}_{\mathsf{C}} \end{array}$$

Upshot If there is a nontrivial  $\{ \Psi_{c} \}_{c} \in \mathbb{Z}(\mathcal{B})^{\times}$  then  $\overline{\Phi}^{\Psi} \neq \overline{\Phi}$ .

Example Let R be a ring, C = R - Mod, and  $u \in R$  an element which is <u>central</u> and a <u>unit</u>. Then for an R-Module M, define

$$\mathcal{Y}_{\mathcal{M}}: \mathcal{M} \longrightarrow \mathcal{M}, \quad \mathcal{Y}_{\mathcal{M}}(\mathbf{x}) = \mathbf{u} \cdot \mathbf{x}.$$

This is R-linear (u is central), and an iso (u is a unit in R), and it is natural (clearly). So as long as  $u \neq 1$ ,  $\{Y_M\}_M$  is a nontrivial unit in the center of R-Mocl, giving a "essentially different" nation  $\mathbb{E}^Y$ .

Example G an abelian group, R=ZG, uEG, u=e.

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