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We have seen the definition of categories and functors, and some simple examples. This lecture enumerates some more examples, leading up to simplicial sets as a motivation for functor categories.

Def A category C is small if ob(C) is a set (not just a class).

<u>Def</u><sup>N</sup> A small category C is <u>ordered</u> if for every pair of objects A, Bthe set C(A, B) contains at most one element.

<u>Lemma</u> If C is ordered then  $(ob(0), \leq)$  is a preorder where we define  $A \leq B$  if and only if C(A, B) is nonempty.

<u>Proof</u> A preorder is a reflexive, transitive relation, so this is clear from the axioms.

Lemma If  $(X, \leq)$  is a preorder then there is a small category  $\mathcal{C}_{X}$  with  $ob(\mathcal{C}_{X}) = X$  and

<u>Def</u><sup>N</sup> A <u>monphism</u> of preorders  $f \colon (X_1, \leq) \longrightarrow (X_2, \leq)$  is a function  $f \colon X_1 \rightarrow X_2$ such that whenever  $x \leq x'$  in  $X_1$ , then  $f(x) \leq f(x')$  in  $X_2$ . This defines the category of preorders <u>Pre</u>.

(2)
<u>Def</u> Let <u>Cat</u> denote the category of small categories and functors.
(the existence of this follows from exercises in Lecture 2).
$\underline{Def}^{N}$ A functor $F: \mathcal{C} \longrightarrow \mathcal{B}$ is
• $f_{aithful}$ if for all pairs $A, B \in ob(\mathcal{O})$ ,
<b>`</b>
$\mathcal{C}(A, \mathcal{B}) \longrightarrow \mathcal{P}(FA, F\mathcal{B}) \tag{2.1}$
is i <u>njective</u> , and
<ul> <li>full if for all pairs A, B ∈ ob(B), (2.1) is surjective, and</li> </ul>
• essentially surjective if for every $D \in ob(P)$ there exists $C \in ob(C)$
and an isomorphism $F(c) \cong D$ .
Fis fully faithful if it is both full and faithful, and an equivalence
if it is fully faithful and essentially surjective.
Lemma There is a fully faithful functor F: Pre -> Cat defined by
$F(X, \leq) = G_X$ on objects and for $f: (X, \leq) \rightarrow (X_2, \leq)$ in free by
$F(f): \mathcal{C}_{x_1} \longrightarrow \mathcal{C}_{x_2}$
F(f)(x) = f(x)
$\sum \left( \frac{1}{1} + \frac{1}{1} +$

 $F(*:x_1 \to x_2) = *:f(x_1) \to f(x_2).$ 

Ex 1 Prove the lemma.

<u>Note</u> This defines an equivalence <u>Pre</u>  $\longrightarrow$  <u>Drd</u> where <u>Ord</u> denotes ordeved categories and functors. <u>Def</u> Let G be a category. A <u>subcategory</u> T of G is a class ob(T) = ob(B)and for each pair  $A, B \in ob(T)$  a subject T(A, B) = C(A, B) s.t. for all  $A \in ob(T)$  we have  $[A \in T(A, A), and T$  is closed under composition in G. Thus T is itself a category, and there is a fuithful functor  $T \longrightarrow C$ .

<u>Def</u><sup>N</sup> For p∈N consider the preorder [P] := {0,1,..., P} with the usual relation ≤. We denote by  $\Delta$  the full subcategoing of <u>Pre</u> with objects {[P]}<sub>p>0</sub>. That is,

$$ob(\mathbb{A}) = \{ [o], [v], \dots \}$$

$$Hom_{\Delta}([P], [q]) = Hom_{Pre}([P], [q])$$
$$= \{f: \{0, 1, ..., p\} \longrightarrow \{0, 1, ..., q\} \mid i \leq j$$
$$implies \ f(i) \leq f(j)\}.$$

(3)

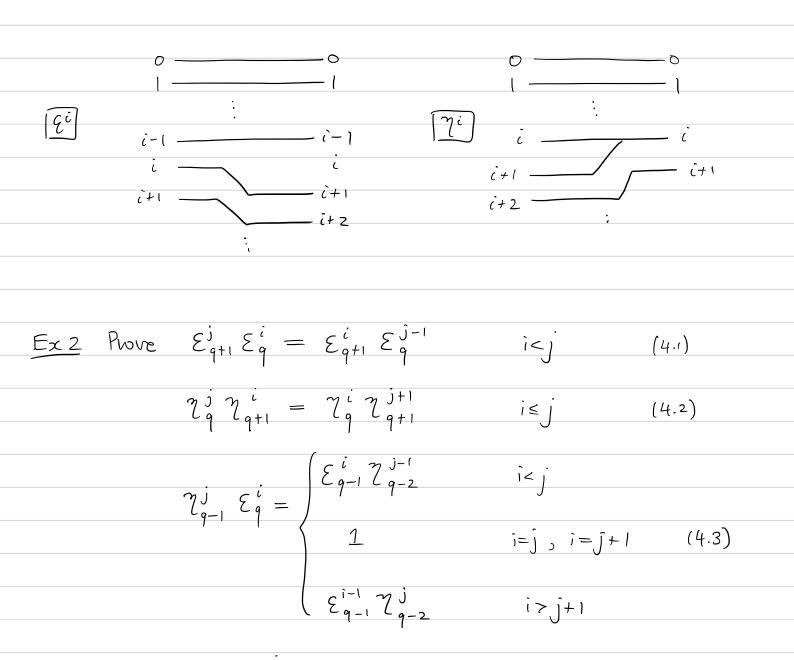
This is quite an interesting and important category, for both topology and homological algebra. It is called (for reasons that will be come clear) the simplex category. By the above, there is a fully faithful functor  $\Delta \longrightarrow Cat$ , if we choose to think of [P] as a category.

<u>Def</u> Given  $0 \le i \le q$  integers, define morphisms in A

$$\mathcal{E}^{i} = \mathcal{E}^{i}_{q} : [q-1] \longrightarrow [q] \qquad \mathcal{N}^{i} = \mathcal{N}^{i}_{q} : [q+1] \longrightarrow [q]$$

$$\mathcal{E}^{i}(a) = \begin{cases} a & a < i \\ a+1 & a \neq i \end{cases} \qquad \mathcal{N}^{i}(a) = \begin{cases} a & a \leq i \\ a-1 & a \neq i \end{cases} \qquad (3.1)$$

where 
$$\mathcal{E}_{q}^{c}$$
 is defined for  $q > 0$ . (3.2)



<u>Note</u> Observe that the E<sup>c</sup> are injective, and since composition in *A* is of functions between sets, this implies E<sup>c</sup> are monomorphisms. Similarly all the N<sup>i</sup> are epimorphisms. In fact something stronger can be read from (4.3).

<u>Def</u><sup>N</sup> In a category Ca morphism  $f: A \rightarrow B$  is a <u>retraction</u> (or <u>splif</u> <u>epimonohism</u>) if there exists  $g: B \rightarrow A$  with  $f \circ g = IB$ . Dually, f is a <u>wretraction</u> (or <u>splif</u> monomorphism) if there exists  $g: B \rightarrow A$  with  $g \circ f = IA$ .

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Lemma Splitepi 
$$\Rightarrow$$
 epi and split mono  $\Rightarrow$  mono.  
Proof Suppose  $f \circ g = la$  and  $uf = vf$ , then  
 $u = ul_B = ufg = vfg = vl_B = v.$   
Similarly for split monos. []  
For  $osiseq$ , and  $\eta q q q t = 1$   
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For  $osiseq$ , and  $\eta q q q t = 1$   
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For  $osiseq$ , and  $\eta q q q t = 1$   
 $similarly for split monos. []
For  $for pointing out this oversight
Lemma In  $\Delta$  every  $E^i$  is split mono and every  $T^i$  is split epi.  
Roof  $ey(4.3)$ ,  $T^i E^i = 1$ , i.e.  $[q] \xrightarrow{e^i} [q+1] \xrightarrow{\gamma} [q] = id_{EQ}$ . []  
Ex. 3 Any composite of epis (resp. split epis) in a category  $G$  is again  
anepi (resp. split epi) and the same for monos  $A$  split monos.  
Theorem Any morphism  $\mu: [q] \rightarrow [q]$  in  $\Delta$  can be curitlen uniquely as  
 $\mu = E^{i_1} \cdots E^{i_s} T^{i_1} \cdots T^{i_s}$   
 $uith p \ge i_1 > i_2 \cdots > i_s \ge 0$  and  $0 \le j_1 < \cdots < j_s < q$  and  
 $q - t + s = p$ . The identify  $\mu = l$  corresponds to  $s = t = 0$ .  
Note  $[q] \xrightarrow{\gamma} (q-1] \xrightarrow{\gamma} \cdots \xrightarrow{\gamma} [q-t] \xrightarrow{\varepsilon} [q-t+1] \longrightarrow \cdots \xrightarrow{\varepsilon} [q-t+s=p]$   
 $responsible for "kemel" of  $\mu$ .$$$$$